

# CS383 Programming Languages

Quiz 7

1. What is the difference between a list and an array?

Array is mutable

2. What do you think are the major advantages of the imperative programming style over the functional style?

It's easier to programmer.

### 3. Which rule is **incorrect**?

- a. 
$$\frac{(M, e) \rightarrow (M', e')}{(M, \text{ref } e) \rightarrow (M', \text{ref } e')} \text{ (E-Ref)}$$
- b. 
$$\frac{l \notin \text{dom}(M)}{(M, \text{ref } v) \rightarrow ((M, l \mapsto v), l)} \text{ (E-RefV)}$$
- c. 
$$\frac{(M, e) \rightarrow (M', e')}{(M, !e) \rightarrow (M', !e')} \text{ (E-DeRef)}$$
- d. 
$$\frac{}{(M, !l) \rightarrow (M, l)} \text{ (E-DeRefLoc)}$$
- $$\frac{}{(M, !l) \rightarrow (M, M(l))} \text{ (E - DeRefLoc)}$$

# References (Operational Semantics, Cont'd)

$$\frac{(M, e) \rightarrow (M', e')}{(M, \text{ref } e) \rightarrow (M', \text{ref } e')} \text{ (E - Ref)}$$

$$\frac{l \notin \text{dom}(M)}{(M, \text{ref } v) \rightarrow ((M, l \mapsto v), l)} \text{ (E - RefV)}$$

$$\frac{(M, e) \rightarrow (M', e')}{(M, !e) \rightarrow (M', !e')} \text{ (E - DeRef)}$$

$$\frac{}{(M, !l) \rightarrow (M, M(l))} \text{ (E - DeRefLoc)}$$

$$\frac{(M, e_1) \rightarrow (M', e'_1)}{(M, e_1 := e_2) \rightarrow (M', e'_1 := e_2)} \text{ (E - Assign1)}$$

$$\frac{(M, e_2) \rightarrow (M', e'_2)}{(M, v_1 := e_2) \rightarrow (M', v_1 := e'_2)} \text{ (E - Assign2)}$$

$$\frac{}{(M, l := v) \rightarrow (M[l \mapsto v], ())} \text{ (E - Assign)}$$

#### 4. Which rule is **incorrect**?

- a. 
$$\frac{(M, e_1) \rightarrow (M', e'_1)}{(M, e_1 := e_2) \rightarrow (M', e'_1 := e_2)} \text{ (E - Assign1)}$$
- b. 
$$\frac{(M, e_2) \rightarrow (M', e'_2)}{(M, v_1 := e_2) \rightarrow (M', v'_1 := e'_2)} \text{ (E - Assign2)}$$
- c. 
$$\frac{}{(M, l := v) \rightarrow (M[l \mapsto v], v)} \text{ (E-Assign)}$$
- d. None of the above
- $$\frac{}{(M, l := v) \rightarrow (M[l \mapsto v], ())} \text{ (E-Assign)}$$

## 5. Which one is incorrect?

a.

$$\frac{\Sigma; G \vdash e_1 : t \quad \Sigma; G \vdash e_2 : t}{\Sigma; G \vdash e_1 := e_2 : \text{unit}} \text{ (T-Assign)}$$

b.

$$\frac{\Sigma; \Gamma \vdash e : t \text{ ref}}{\Sigma; \Gamma \vdash !e : t} \text{ (T - Deref)}$$

c.

$$\frac{\Sigma; \Gamma \vdash e : t}{\Sigma; \Gamma \vdash \text{ref } e : t \text{ ref}} \text{ (T - Ref)}$$

d.

$$\frac{\Sigma(l) = t}{\Sigma; \Gamma \vdash l : t \text{ ref}} \text{ (T - Loc)}$$

$$\boxed{\frac{\Sigma; \Gamma \vdash e_1 : t \text{ ref} \quad \Sigma; \Gamma \vdash e_2 : t}{\Sigma; \Gamma \vdash e_1 := e_2 : \text{unit}}} \text{ (T - Assign)}$$

# REFERENCES (TYPING)

- We define the typing relation for memory store as  $\Sigma$  (or  $S_i$ ):

$\Sigma ::= . \mid \Sigma, l : t$  ( $t$  is the type of value stored at  $l$ )

- Our new typing judgment:

$$\Sigma; \Gamma \vdash e : t$$

- Types:  $t ::= .. \mid \text{unit} \mid t \text{ ref}$

$$\frac{}{\Sigma; \Gamma |- x : \Gamma(x)} \text{ (T - Var)}$$

$$\frac{\Sigma; \Gamma, x : t_1 |- e : t_2}{\Sigma; \Gamma |- \lambda x : t_1. e : t_1 \rightarrow t_2} \text{ (T - Abs)}$$

$$\frac{\Sigma; \Gamma |- e_1 : t_1 \rightarrow t_2 \quad \Sigma; \Gamma |- e_2 : t_1}{\Sigma; \Gamma |- e_1 \ e_2 : t_2} \text{ (T - App)}$$

$$\frac{}{\Sigma; \Gamma |- () : \text{unit}} \text{ (T - Unit)}$$

$$\frac{\Sigma(l) = t}{\Sigma; \Gamma |- l : t \text{ ref}} \text{ (T - Loc)}$$

$$\frac{\Sigma; \Gamma |- e : t}{\Sigma; \Gamma |- \text{ref } e : t \text{ ref}} \text{ (T - Ref)}$$

$$\frac{\Sigma; \Gamma |- e : t \text{ ref}}{\Sigma; \Gamma |- !e : t} \text{ (T - Deref)}$$

$$\frac{\Sigma; \Gamma |- e_1 : t \text{ ref} \quad \Sigma; \Gamma |- e_2 : t}{\Sigma; \Gamma |- e_1 := e_2 : \text{unit}} \text{ (T - Assign)}$$

6. In an expression  $e_1; e_2$ , if  $e_1$  doesn't evaluate to  $()$ , what will happen?

Raise error(get stuck)

## 7. Which one is incorrect?

- a. 
$$\frac{(M, e_1) \rightarrow (M', e'_1)}{(M, e_1; e_2) \rightarrow (M', e'_1; e_2)} \text{ (E-Seq1)}$$
- b. 
$$\frac{}{(M, () ; e) \rightarrow (M, e)} \text{ (E-Seq2)}$$
- c. 
$$\frac{\Sigma ; \Gamma |- e_1 : \text{unit} \quad \Sigma ; \Gamma |- e_2 : t}{\Sigma ; \Gamma |- e_1 ; e_2 : t} \text{ (T-Var)}$$
- d. All of them are correct

8. Write down the evaluation steps of the following expression:

```
let y = ref 1 in  
  let f = \x. x + !y in  
    f (!y)
```

| 1, let y = | in  
 let f = \x. x + !y in  
 f (!y)

| 1, let f = \x. x + !! in  
 f (!!)

| 1, \x.x+!! (!!)

| 1, \x.x+!! 1

| 1, 1+!!

| 1, 1+1

| 1, 2

## 9. Which one is incorrect ?

- a.  $\frac{\text{try } v \text{ with } e \rightarrow v}{(E - \text{TryV})}$
- b.  $\frac{\text{try error with } e \rightarrow \text{error}}{(E - \text{TryError})}$   $\frac{\text{try error with } e \rightarrow e}{(E - \text{TryError})}$
- c.  $\frac{e_1 \rightarrow e_1'}{\text{try } e_1 \text{ with } e_2 \rightarrow \text{try } e_1' \text{ with } e_2} (E - \text{Try})$
- d.  $\frac{\Gamma |- e_1 : t \quad \Gamma |- e_2 : t}{\Gamma |- \text{try } e_1 \text{ with } e_2 : t} (T - \text{Try})$

# HANDLING EXCEPTION

- **Syntax:**

$e ::= \dots$   
|  $\text{try } e_1 \text{ with } e_2$  (trap errors)

- **Evaluation:**

$$\frac{}{\text{try } v \text{ with } e \rightarrow v} \text{ (E-TryV)}$$

$$\frac{}{\text{try error with } e \rightarrow e} \text{ (E-TryError)}$$

$$\frac{e_1 \rightarrow e_1'}{\text{try } e_1 \text{ with } e_2 \rightarrow \text{try } e_1' \text{ with } e_2} \text{ (E-Try)}$$

- **Typing:**

$$\frac{\Gamma |- e_1 : t \quad \Gamma |- e_2 : t}{\Gamma |- \text{try } e_1 \text{ with } e_2 : t} \text{ (T-Try)}$$

## 10. What is the problem with the following rule?

$$\frac{}{G \vdash \text{error} : t} \quad (\text{T-Error})$$

t can be any type:

$(\lambda x:\text{bool} . x)$  error      error: bool

$(\lambda x:\text{bool} . x) (\text{error} \text{ true})$  error: bool  $\rightarrow$  bool

This breaks the uniqueness lemma

subtyping, or polymorphic types