

6.1

Suppose j is the number of seen symbols in the stream. Currently we are seeing a_{j+1} . We denote the sampling symbol as a_i . We maintain j and a_i . a_i is updated to a_{j+1} with probability $\frac{1}{j+1}$, and with probability $1 - \frac{1}{j+1}$ remains the old value. If we keep a_i as our selection, then it will have been selected with probability $\left(1 - \frac{1}{j+1}\right)^j = \frac{1}{j+1}$, which is the correct probability for selecting a_i from the stream. Therefore, the space we have to maintain is $O(\log m + \log n)$.

6.3

see solution to 6.1.

each symbol is now a word.

6.8

$$H = \{h(x) = ax \bmod M\}$$

6.9

(a) No

Since the set $\{h_{ab}\}$ can be determined by any two of equations $h_{ab}(x) = u$, $h_{ab}(y) = v$, $h_{ab}(z) = w$, this set is not of hash functions 3-universal.

(b) $H = \{h_{abc}(x) = ax^2 + bx + c \bmod p \mid 0 \leq a, b, c < p\}$

With $h_{abc}(x) = u$, $h_{abc}(y) = v$, $h_{abc}(z) = w$, we can get

$$\begin{pmatrix} x^2 & x & 1 \\ y^2 & y & 1 \\ z^2 & z & 1 \end{pmatrix} \begin{pmatrix} a \\ b \\ c \end{pmatrix} = \begin{pmatrix} u \\ v \\ w \end{pmatrix} \pmod{p}$$

This equation has solution $\begin{pmatrix} a \\ b \\ c \end{pmatrix} = \begin{pmatrix} x^2 & x & 1 \\ y^2 & y & 1 \\ z^2 & z & 1 \end{pmatrix}^{-1} \begin{pmatrix} u \\ v \\ w \end{pmatrix}$, because $x \neq y \neq z$.

So, there is a unique $\begin{pmatrix} a \\ b \\ c \end{pmatrix}$ that meet the criteria. Hence

$$\begin{aligned} & \text{Prob}(h_{abc}(x) = u, h_{abc}(y) = v, h_{abc}(z) = w) \\ &= \text{Prob}\left(\text{solution} \begin{pmatrix} a \\ b \\ c \end{pmatrix} = \begin{pmatrix} x^2 & x & 1 \\ y^2 & y & 1 \\ z^2 & z & 1 \end{pmatrix}^{-1} \begin{pmatrix} u \\ v \\ w \end{pmatrix}\right) \\ &= \frac{1}{p} * \frac{1}{p} * \frac{1}{p} \end{aligned}$$

6.10

Take $k=3$, Let $H = \{(0,0,0), (0,1,1), (0,2,2), (1,0,1), (1,1,2), (1,2,0), (2,0,2), (2,1,0), (2,2,1)\}$