

CS383 Programming Languages

Quiz 3

1. Which one of the following is **not** a possible lambda expression?

a. Variable: x

b. Condition: $\text{if } e_1 \text{ then } e_2 \text{ else } e_3$

c. Abstraction: $\lambda x . e$

d. Application: $e_1 e_2$

2. Which one of the following statements is correct?

- a. a name and the object it denote are the same thing.
- b. an object can have only one name.
- c. a name can denote different objects at different times.
- d. λ -calculus uses dynamic binding

3. What are the free variables in the following lambda expression?

$x (\lambda y \cdot y z x) (\lambda m \cdot \lambda n \cdot l m n)$

a. x, z, l

b. x, y, z, l, m, n

c. x, z, n

d. y, l, m

*4. For the following substitutions, which is **incorrect**?

a. $x[y/x] = y$

b. $\lambda x \cdot z w[y/x] = \lambda y \cdot z w$

c. $(\lambda x \cdot (x z))[v/x] = (\lambda y \cdot (v y))[z/y]$

d. $x z w[y/x] = y z w$

Substitution:

- a. $x [[e/x]] = ?$
- b. $y [[e/x]] = ?$ (if $y \neq x$)
- c. $(\lambda x.e1) [[e/x]] = ?$
- d. $(\lambda y.e1) [[e/x]] = ?$ (if $y \neq x$)

6. What's the result of the following lambda expression, under **full beta-reduction**?

$$(\lambda x \cdot x)((\lambda x \cdot x)(\lambda z \cdot (\lambda x \cdot x)z))$$

Full beta-reduction: any redex

a. $\lambda z \cdot z$

b. x

c. z

d. $\lambda z \cdot \lambda z \cdot z z$

5. Which one of the following is **different** from the other three after **call by name** evaluation?

Call-by-name: leftmost, outermost redex first,
NO reduction inside lambda abstractions

a. $\lambda x. x x$

b. $(\lambda x. x x) (\lambda y. y y)$

c. $(\lambda y. \lambda x. y x) (\lambda x. x x) (\lambda y. y y)$

d. $(\lambda x. (\lambda x. x x) x) (\lambda x. x x)$

7. What is the **first step** of
 $(\lambda y. (\lambda x. x) y)$ $((\lambda u. u) (\lambda v. v))$
under call-by-**name** evaluation?

Call-by-name: leftmost, outermost redex first,
NO reduction inside lambda abstractions

a. $(\lambda y. y) ((\lambda u. u) (\lambda v. v))$

b. $(\lambda x. x) ((\lambda u. u) (\lambda v. v))$

c. $(\lambda y. (\lambda x. x) y) (\lambda v. v)$

d. $(\lambda y. (\lambda x. x) y) (\lambda u. u)$

*8. What is the **first step** of
 $(\lambda y. (\lambda x. x) y) \underline{((\lambda u. u) (\lambda v. v))}$
under **call-by-value** evaluation?

call-by-value: only outermost redex, whose RHS must be a value (λ abstraction), no reduction inside abstraction

a. $(\lambda y. y) ((\lambda u. u) (\lambda v. v))$

b. $(\lambda x. x) ((\lambda u. u) (\lambda v. v))$

c. $(\lambda y. (\lambda x. x) y) (\lambda v. v)$

d. $(\lambda y. (\lambda x. x) y) (\lambda u. u)$

9. Application associate to the _____?

*e.g. $M N L = (M N) L$ – associate to the left
or $M(N L)$ – associate to the right*

a. Left

b. Right

*10. What's equivalent to $\lambda y. y z \lambda x. x y z$?

- a. $\lambda y. \underline{(y z)} \lambda x. \underline{((x y) z)}$ - extends as far as possible to the right?
- b. $\lambda y. (y (z \lambda x. \underline{(x (y z))}))$ - association to the left?
- c. $(\lambda y. (y z)) \lambda x. \underline{((x y) z)}$ - extends as far as possible to the right?
- d. $\lambda y. ((y z) \lambda x. ((x y) z))$